Automating Custom-Precision Function Evaluation for Embedded Processors

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Talk outline

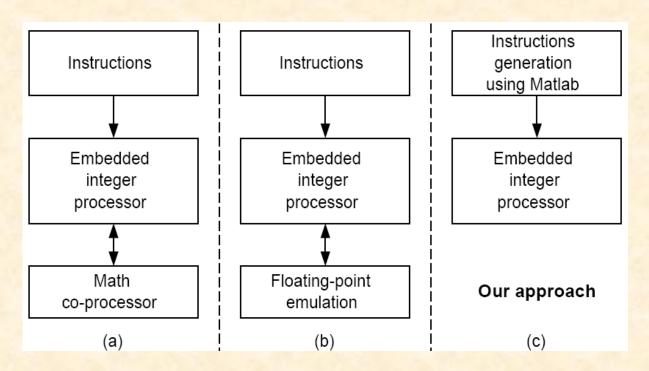
- 1. achievements
- 2. motivation
- 3. function evaluations
- 4. design tool flow
- 5. error analysis
- 6. performance evaluation
- 7. future work
- 8. summary

1. Achievements

- customizable library for floating-point function evaluation based on input integer instruction set
- automatic code generation using high-level Matlab model, and optimization for customizing precisions
- evaluation of this method with two elementary functions and Xilinx embedded design kit
- automating the selection of approximation method, polynomial degree for a given function, accuracy requirement and execution time

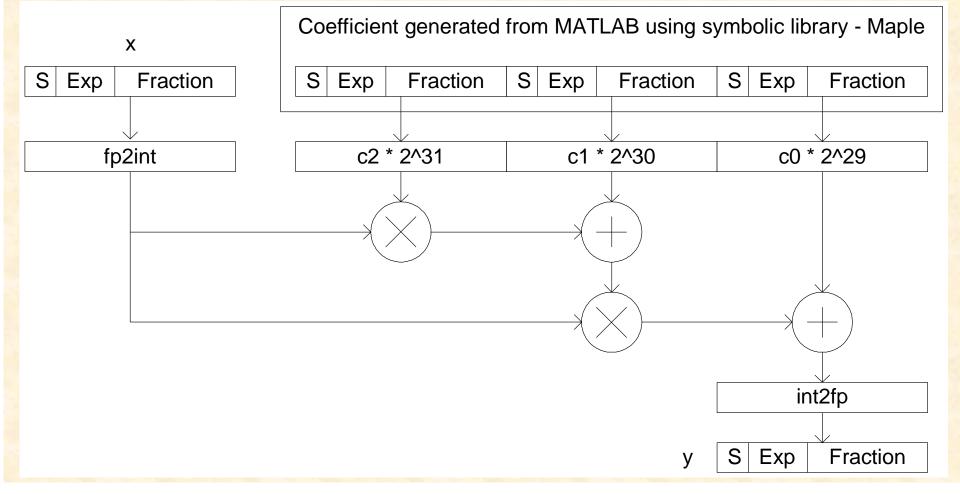
2. Motivation

- embedded systems are usually space and time critical, a dedicated coprocessor and a larger memory for instruction are infeasible
- previous work on math co-processor and floating point emulation
- automated code generation for mathematical function library targeting customizable precision (depending on the error requirements)



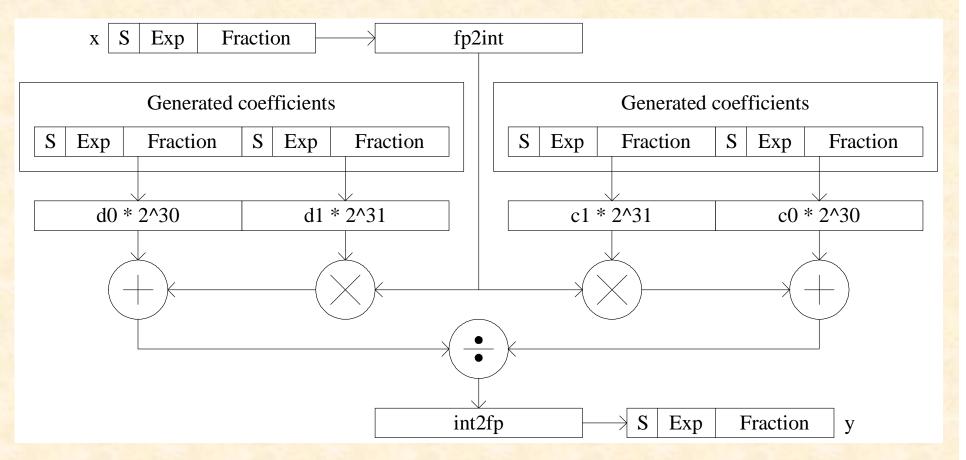
3. Function approximation

Polynomial method: approximation with a single polynomial



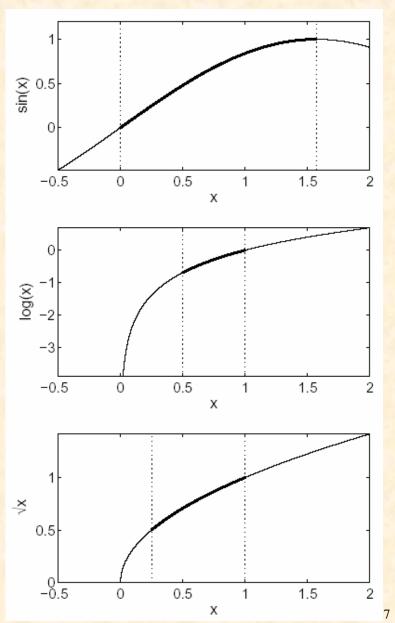
Rational approximation

Rational method: with two polynomials (same degree)



Range Reduction

- f(x) where x=[a,b]
 - (1) range reducing x to a more convenient interval y=[a',b']
 - (2) function approximation on the reduced interval
 - (3) range reconstruction:
 expanding the result back
 to the original result range

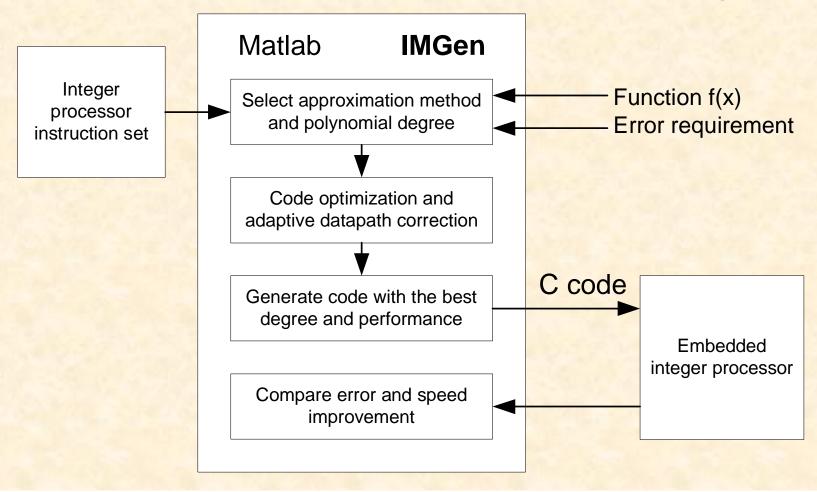


Example: Evaluating log(x)

```
Evaluating f(x) = \log(x)
                                             input x
// Range Reduction
input.sng_as_flt = x;
exp = input.sng_as_fld.exp - 126;
ix = fp2int(input); // y = ix;
// Evaluation Method
// f(y) where y = [0.5, 1)
// e.g. degree-3 polynomial
f1 = ((c_3 \times y + c_2) \times y + c_1) \times y + c_0;
// Range Reconstruction
s1 = range(exp); // find \exp \times \log(2)
f1 = (f1 \gg overflow) + s1;
                                          Adjust the
output = int2fp(f1);
                                          output exponent
output.sng_as_fld.exp += overflow;
```

4. Design tool flow using Matlab

- technology-independent flow
 - use the embedded PowerPC as an example



IMGen – 3 steps

- automation:
 - user error requirement → function evaluation implementation
 - select rational / polynomial approximation
 - select rational / polynomial degree
 - select 32-bit / 64-bit datapath
- generation: using custom precision code
 - Add / multiply / divide / shift operations
- optimization: e.g. loop unrolling techniques

Code optimization

code generation optimization example

```
long ic[2] = \{1488522235, -1456492463\};
 #define CORRECTION if (j==0) s3 = s3 << 1;
unoptimized code
 ix = fp2int(input);
 iy = ic[0];
 for (j=0; j<degree; j++){</pre>
    mhw(ix, iy, s3);
    CORRECTION
    iy = s3 + ic[j+1];
using loop-unrolling technique
 ix = fp2int(input);
 mhw(ix, ic[0], s3);
```

s3 = s3 << 1;

iy = s3 + ic[1];

generated coefficient

adaptive datapath correction

Floating-point - fixed-point conversion

- input and output are both floating-point format
- internal computation is transparent to users

```
C-code: Data structure for input/output
typedef struct sng_flds {
unsigned sgn : 1; // 0x8000 0000
unsigned exp: 8; // 0x7F80 0000(bias 127)
unsigned val: 23; // 0x007F FFFF
} SNG_FLD;
C-code: fp2int - floating-point to integer
output = input.val << 8;</pre>
output += 0x80000000;
output = output >> 1;
```

5. Error analysis

- approximation error (error of approximating a function, e.g. using minimax)
- quantization error induced by: (1) the multiply-add datapath, (2) range reconstruction (function dependent)
- rational approximation has a much lower approximation error

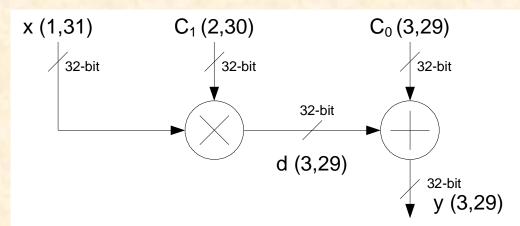
$$E_{total} = E_{approximation} + E_{quantization} \\ E_{quantization} = E_{poly_approximation} + E_{range_reconstruction}$$

Degree	Polynomial approx.	Rational approx.
1	0.02983005	0.0008607941
2	0.00342398	0.0000017146
3	0.00044161	0.0000000032

Error analysis example

- quantization error analysis of degreeone log(x)
- x(1,31) = 31 fraction bits
- E_d = error accumulated at signal d
- higher degree

 higher error



$$E_y = E_{poly_approximation}$$

$$E_y = E_{accum} + E_{quantization_y}$$

$$E_y = (E_d + E_{C_0}) + 2^{-29}$$

$$E_d = E_{C_1} + 2^{-29}$$

$$E_{C_1} = 2^{-31}$$

$$E_{C_0} = 2^{-30}$$

Therefore,

$$E_y = 2^{-31} + 2^{-29} + 2^{-30} + 2^{-29}$$

 $E_y = 5.12227 \times 10^{-9}$

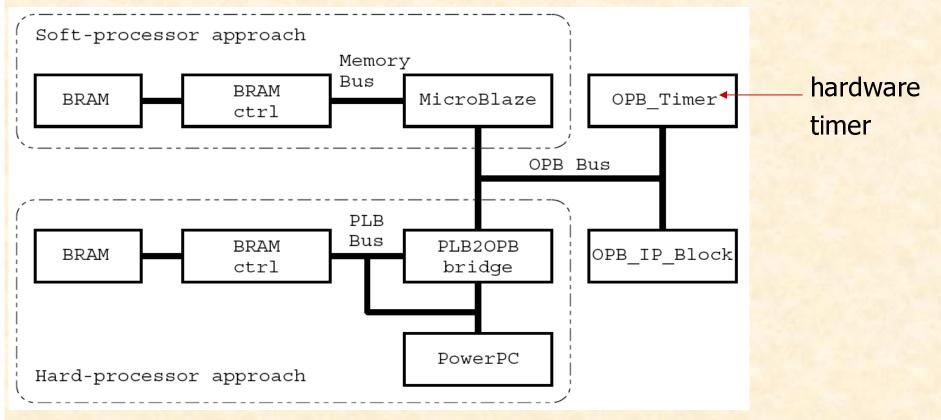
System automation

 input / output via Matlab, remote execution on the embedded system board

```
User input >> genlib('log', 0.01)
Phase 1: Maple command generates polynomial coefficients
Phase 2: Static error analysis calculates quantization error
Phase 3: Select polynomial/rational approximation
Phase 4: Select 32-bit/64-bit implementation
Phase 5: Generate embedded C code
         and execute in embedded integer processor
Phase 6: Output performance data and statistical error
       text data bss
                         dec hex
                                    filename
      44232 4296 48 48576 bdc0 TestApp/executable.elf
      cycle count for the Xilinx math library: 63335
      cycle count for the bus overhead:
                                              60
      cycle count for the IMGen library:
                                              618
      average speedup: 1.13e+002
      maximum error : 0.0034241
      IMGen is generated and tested in 4.904e+001 seconds
```

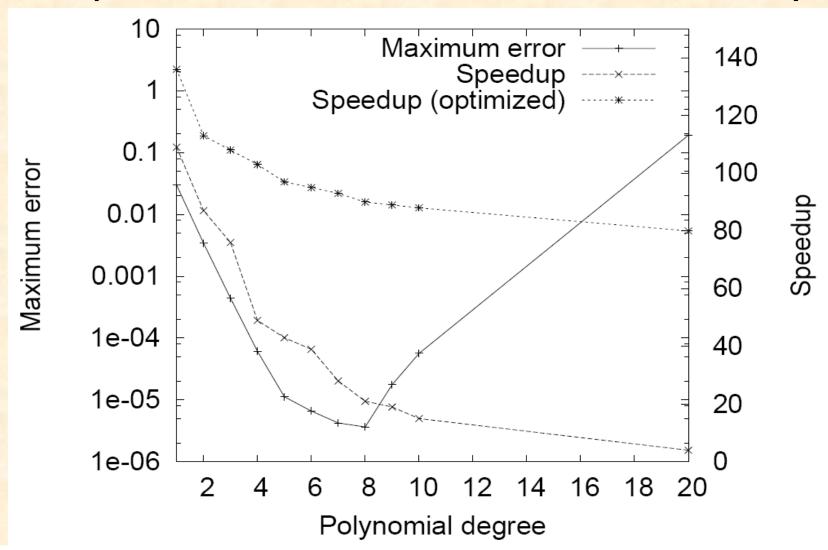
Embedded system under test

- use Xilinx ML310 system, XC2VP30 device, with two embedded PowerPC chips
- can target Xilinx MicroBlaze soft integer processor



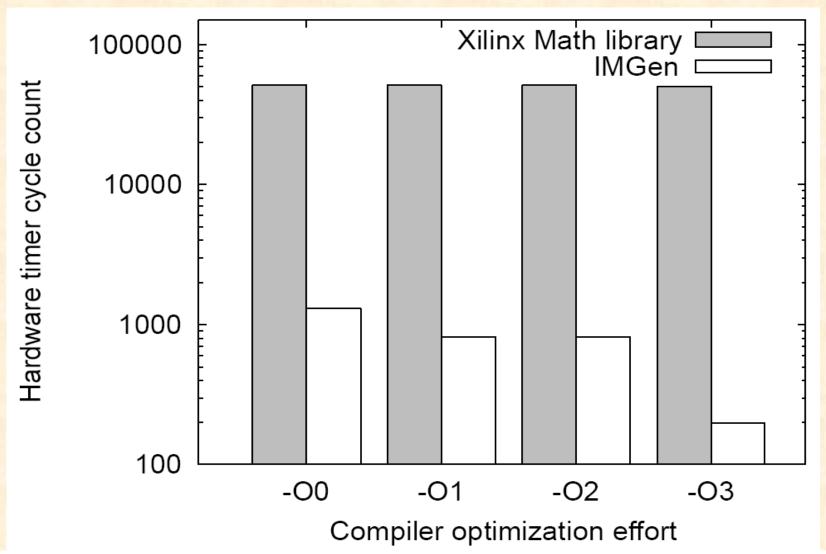
6. Performance evaluation

compare with Xilinx emulated math library



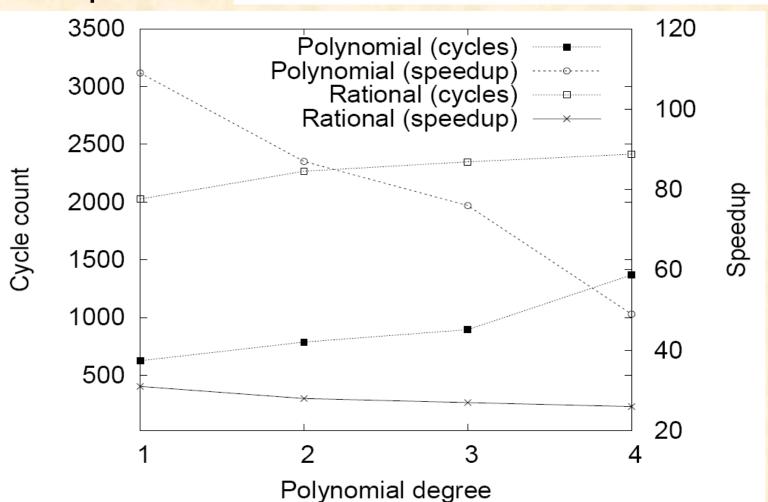
Compile time optimization

study the effect of compiler optimization



Polynomial vs. rational

• we measure the bus latency, measure an accurate speedup factor $Speedup = (T_{math_library} - T_{overhead})/(T_{IMGen} - T_{overhead})$



Performance comparisons

	32-bit	32-bit (opt.)	64-bit
Xilinx Math (cycles)	63759	63699	64370
Bus latency (cycles)	60	60	60
IMGen (cycles)	1369	672	1921
Speedup factor	48x	103x	34x
Measured error	0.00005886	0.00005920	0.00000159

tradeoff between speed and accuracy

	$\log\left(x\right)$	\sqrt{x}
Xilinx Math (cycles)	62725	9159
Bus latency (cycles)	60	60
IMGen (cycles)	1696	467
Speedup factor	38x	22x
Measured error	0.000004313	0.0008375

7. Future work

- code generation for more integer processors
- comparison with floating-point coprocessor
- use better range reduction technique for software implementation
- use run-time reconfiguration to configure soft-processors such as MicroBlaze

8. Summary

- customizable library for floating-point function evaluation based on input integer instruction set
- automatic code generation using high-level Matlab model, and optimization for customizing precisions
- evaluation of this method with two elementary functions and Xilinx embedded design kit
- automating the selection of approximation method, polynomial degree for a given function, accuracy requirement and execution time
- embedded code generator
 - cope with speed/code-size/error trade-off